Software Upgrades for Distributed Systems

> Sameer Ajmani Google

Barbara Liskov Liuba Shrira MIT Brandeis

Internet Services

- v Are long-lived, robust
- v Run on many machines
- Must be continuously available
- Have persistent state
- Face ever-changing requirements
- Require software upgrades to
 - Fix bugs
 - Improve performance
 - v Add/change/remove features

Upgrade Requirements

- Automatic, Controlled Deployment
 - Ensure continuous availability
 - Test new software on a few nodes
 - Upgrade servers before clients
- **v** Mixed mode operation



Outline

- v System & Upgrade Model
- v Specifying Upgrades
- v Implementation Models

System Model

- v A node is an object of class C
- Different nodes may run different classes
- v Nodes communicate via RPCs

Upgrade Model

- $\label{eq:class} \begin{array}{l} \nu \quad A \mbox{ class upgrade replaces an old class,} \\ C_{\mbox{old}}, \mbox{ with a new one, } C_{\mbox{new}} \end{array}$
 - Implements types T_{old}, T_{new}
 - May be compatible or incompatible
- v An upgrade is a set of class upgrades

Supporting Mixed Mode

- Each node handles calls to past and future versions of itself
- Adding support for new versions must be fast
- Removing support for old versions should be easy

Simulation Objects

 Each node handles calls to past and future versions of itself

- Future SOs simulate future behavior
- Past SOs simulate past behavior
- Adding/removing an SO does not require a restart



 Each node handles calls to past and future versions of itself



SOs are only required for certain upgrades

Specifying Upgrades

Specifying Upgrades



- Must behave like a single object
 - v Even when upgrades are incompatible
- Upgrade specification must define this
 - v Goal: no surprises for clients!
 - E.g., changing permissions to ACLs

Constraints on Specifications

v Type requirement

 A call of version V behaves according to the specification of type T_V

Constraints on Specifications

v Sequence requirement

- Each event must reflect all earlier ones despite:
 - Client upgrades
 - Server upgrades
 - v Version introduced
 - v Version retired

Example

- v ColorSet & FlavorSet
 - Incompatible upgrade
- ColorSet methods: insertColor(x, c), getColor(x), ...
 - v E.g., { (1, red), (2, blue), (3, red) }
- FlavorSet methods: insertFlavor(x,f), getFlavor(x), ...

Specifications: Invariant

Invariant I relates the object states $I(O_{old}, O_{new})$



I: $\{x | < x, c > in O_{CS}\} = \{x | < x, f > in O_{FS}\}$

Specifications: Mapping Function

Mapping function MF defines initial state $O_{new} = MF(O_{old})$ s.t. I(O_{old} , O_{new})



 $O_{FS} = MF(O_{CS}) = \{<x,grape>| <x,c> in O_{CS}\}$

Relating Behavior



Only for mutators

Relating Behavior



Shadow methods relate behavior Told.m ◊ Tnew.\$m Tnew.p ◊Told.\$p

Shadow Method Specification

void ColorSet.\$insertFlavor(x, f) Effects: no <x,c> in this_{pre} \Rightarrow this_{post} = this_{pre} U {<x,blue>}

 Also ColorSet.\$delete, FlavorSet.\$insertColor, FlavorSet.\$delete

The Compound Type

- I, MF, and the shadow methods define a compound type T_{old&new}
 - All the methods, with extended specs for mutators
- We would like:
 - v T_{old&new} is a subtype of T_{old}, T_{new}
- v When this doesn't work:
 - Weaken invariant I
 - Upgrade scheduling
 - Disallow methods

Implementation Models

Direct Model

SO handles calls only to its own version
 And delegates down the chain



Problems with Direct Model

- v Poor expressive power
 - E.g., FlavorSet SO doesn't know about
 C.insertColor call
- v Synchronization



Interceptor Model

Newest SO gets all calls And delegates down the chain



Interceptor Model

After first (incompatible) upgrade is installed



Interceptor Model

After second (possibly compatible) upgrade is installed



Interceptor Model Evaluation

- v Excellent expressive power
- v Future and past SO must do more
 - v Can reuse code and delegate



Prototype Implementation: Upstart

- v C++ and Sun RPC
- v Intercepts socket(), read(), write()
- Imposes minimal overhead

Summary

- Upstart is the first complete approach
 Allows mixed mode operation
- The first definition of what must be specified for incompatible upgrades
- A powerful and useful implementation model
- A prototype implementation

Software Upgrades for Distributed Systems

Sameer Ajmani

Google Barbara Liskov Liuba Shrira MIT Brandeis

Disallowing

- Constraint: never disallow methods of the current object
- v Future SO may disallow T_{new} methods
- v Past SO may disallow T_{old} methods
- v In either case, disallow
 - Mutators whose shadows are problem
 - v Observers that expose problems

Some shadows cannot be implemented via delegation

- Disallow methods that have unimplementable shadows
- Add shadow method to delegate via dynamic updating
 - Allowed iff T + shadow is a subtype of T
 - E.g., can't add delete() to GrowSet
- Implement shadow method in interceptor
 - Impacts transform function
 - Won't work for past SOs because of retirement

What does Google do?

- Extensible protocols
 - Assume defaults for missing fields
 - v Ignore unexpected fields
- v Round-robin upgrades among replicas
- v Datacenter-by-datacenter

END OF SLIDES

 The remaining slides are leftover from previous talks and may contain stale information

Code Execution

Call contains version number

Called node dispatches



Upstart

- A system that supports upgrades
- v And a methodology
- v Joint work with
 - Barbara Liskov
 - Liuba Shrira

Class Upgrade

- New and old classes C_{new}, C_{old}
 Implement T_{new}, T_{old}
- v Scheduling function SF
- v Transform function TF
- $_{v}$ Simulation classes S_{new} , S_{old}
- Might be incompatible
 T_{new} is not a subtype of T_{old}

Class Upgrade

- Replaces an old class, C_{old} with a new one, C_{new}
- Every node running old class will switch to new class eventually
- Upgrade is a set of class upgrades



- Upgrader enters new upgrade at UDB
- Defines a new version



- Nodes query the UDB periodically
- Version numbers flow on all messages

Executing an Upgrade

- v If upgrade affects the node
 - Runs the SF
 - And simulates the future
 - Shuts down, restarts, runs TF
 - Starts up "normally"
 - And simulates the past

Disallowing Example

- v GrowSet ◊ IntSet
- v For the future SO:
 - v Disallow IntSet.delete
- v For the past SO:
 - v Disallow GrowSet.isIn
- $_{\rm v}~{\rm T}_{\rm old\&new}$ becomes <T_{\rm future} , T_{past}>

What Disallowing Provides

- $_{\rm v}~{\rm T}_{\rm future}$ is a subtype of ${\rm T}_{\rm old}$
- v And it implements T_{new}
- $_{v}$ T_{past} is a subtype of T_{new}
- $_{\rm v}$ And it implements T_{old}

Transform Functions

- Implement the identity map
- May need to use future SO, create past SO
- v Must be restartable
- v Cannot make remote calls

Scheduling Functions

- $_{v}$ Can consult the UDB
- v Examples:
 - Rolling upgrade
 - Big flip
 - Fast reboot

Implementing Upgrades

- v Need to provide SOs, TF, SF
- For the SOs, need an implementation model

Summary of Specifications

- Specification defines the compound type T_{old&new}
 - v I, MF, and the shadows
- If the compound type isn't a subtype, disallow

Specifications: Shadow Methods

v Shadow methods relate behavior



e.g., FlavorSet.\$insertColor

Talk Outline

- v Upgrade requirements
- Upstart overview
- v Specifying upgrades
- Implementing upgrades

Requirement: Generality

- Support for arbitrary changes
- Incompatible upgrades
 - V Old features are no longer supported

Requirement: Continuous Availability

- Service is required 24/7
- v Even when upgrading
- v Therefore systems upgrade gradually
- Implies mixed-mode operation

Requirement: Controlled Deployment

- Systems upgrade gradually
 But with control
- Manual control is impractical
 - An automatic system
 - v But upgrader needs control

Requirement: Persistence

- Systems store important state for users
 - It cannot be lost
 - v But may need to be transformed

Requirement: Ease of Use

- v Avoiding feature creep helps
- Upgrader needs to understand only a few recent versions